

# Rate Based Congestion Control for the Internet

(work in progress)

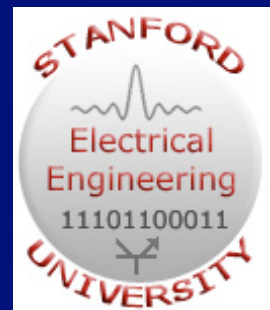


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# Internet Congestion Control

Goals of a congestion control scheme:

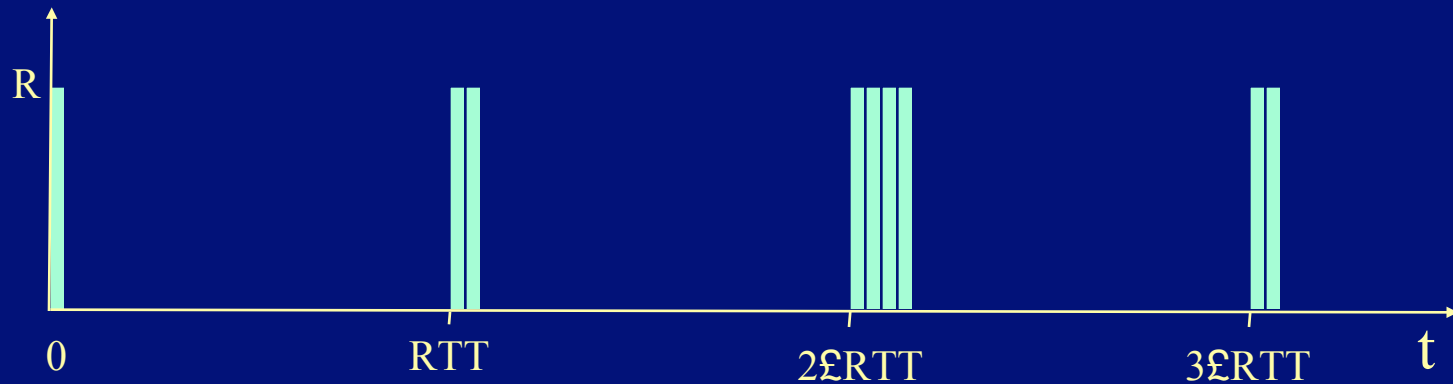
- Limit flow rates to **avoid “congestion”**
- Use the network resources efficiently to **minimize flow durations**
- Ensure **fairness** in resource allocation among flows

# Problem Statement

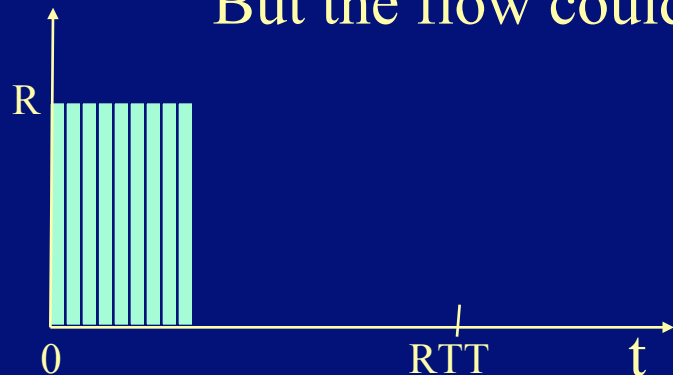
- Finding congestion control scheme for **high bandwidth Internet**
  - Link capacities will continue to increase
  - Flow sizes remain relatively constant
  - More and more flows could finish within a round-trip time (RTT)
- Current congestion schemes:
  - Feedback based
  - Force flows to last multiple RTTs
- **Short flows: flow size/link capacity  $\ll$  RTT**

# Example: A Short Flow

## TCP Congestion Control



But the flow could finish within one RTT!



# Our Scheme for Congestion Control

- Open loop based:
  - No feedback
  - Flow rate determined at start by interaction between routers and end-host
- Flow/user centric instead of packet centric
- Fair amongst flows: flows have same rate
- Minimize response times

# Assumptions

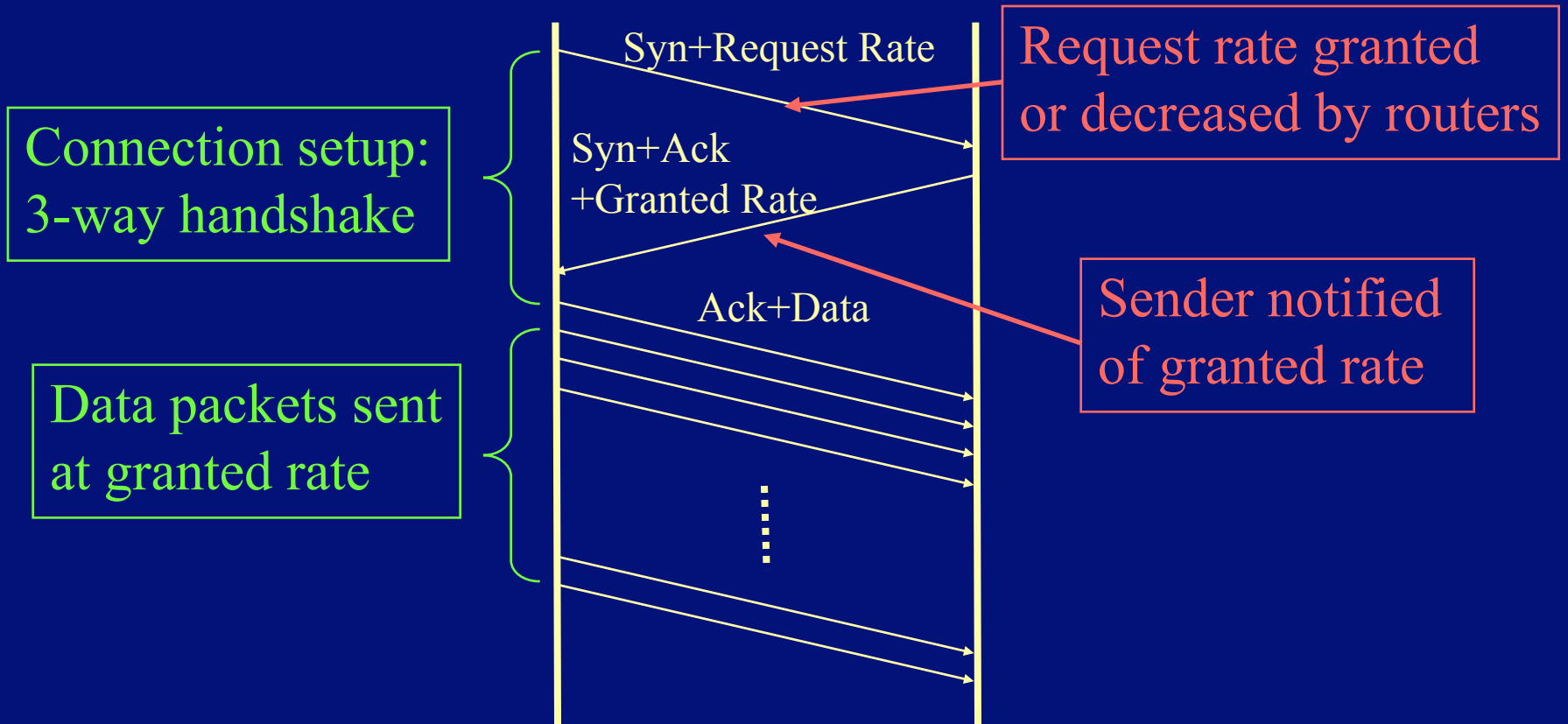
- Network consists of short flows only
- Flows arrive according to a Poisson Process
- Flow sizes are independent and identically distributed

# Goals of Our Scheme

- With high flow rate and finite buffer
  - Probability that router is overloaded is large
  - High chance of queue overflow
- We want our scheme to:
  - Alleviate these undesirable effects
  - Provide low flow response time
  - At the same time, be fair amongst flows

# Rate Control Protocol (RCP)

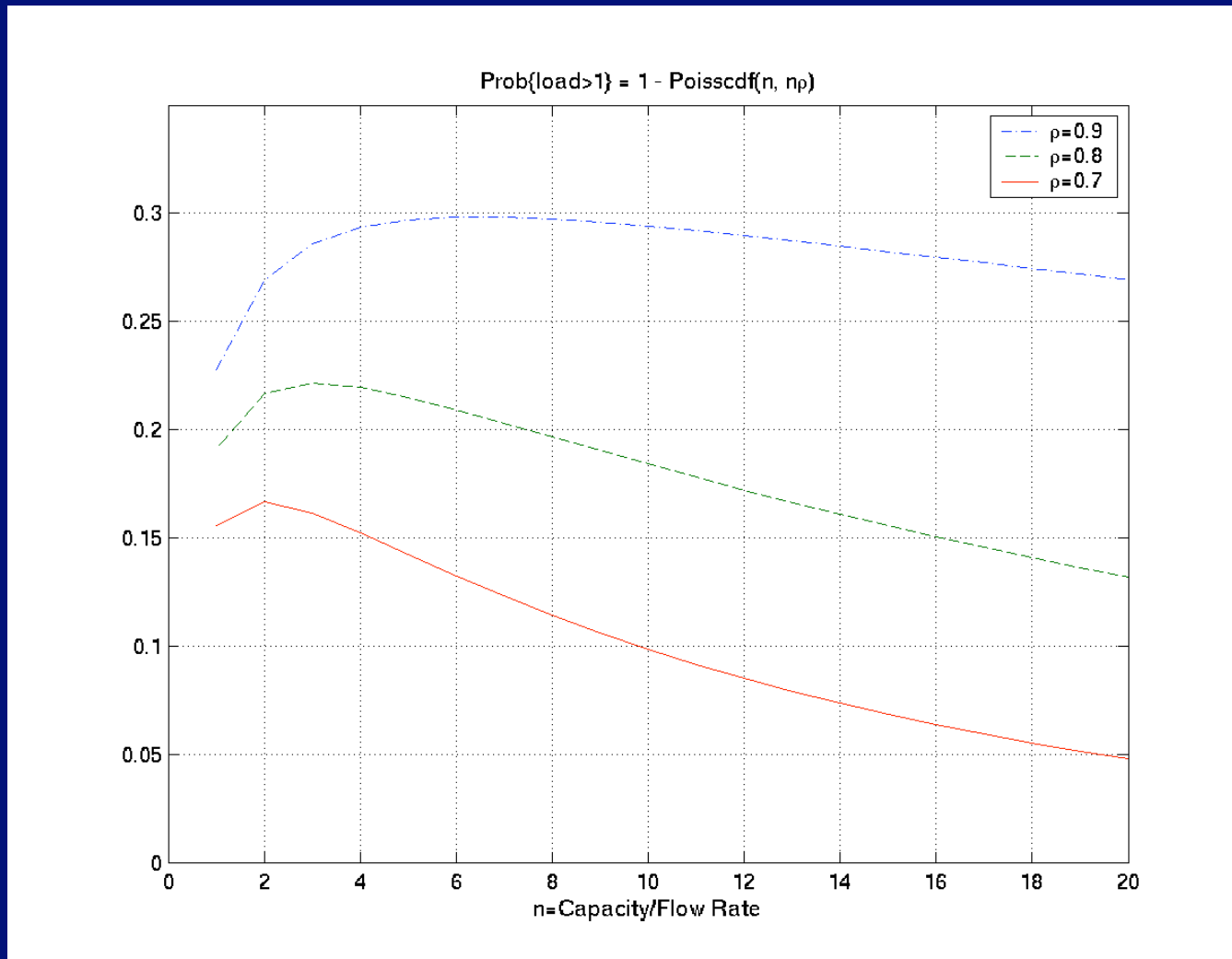
Sender Receiver



# Rate Control Protocol: Router

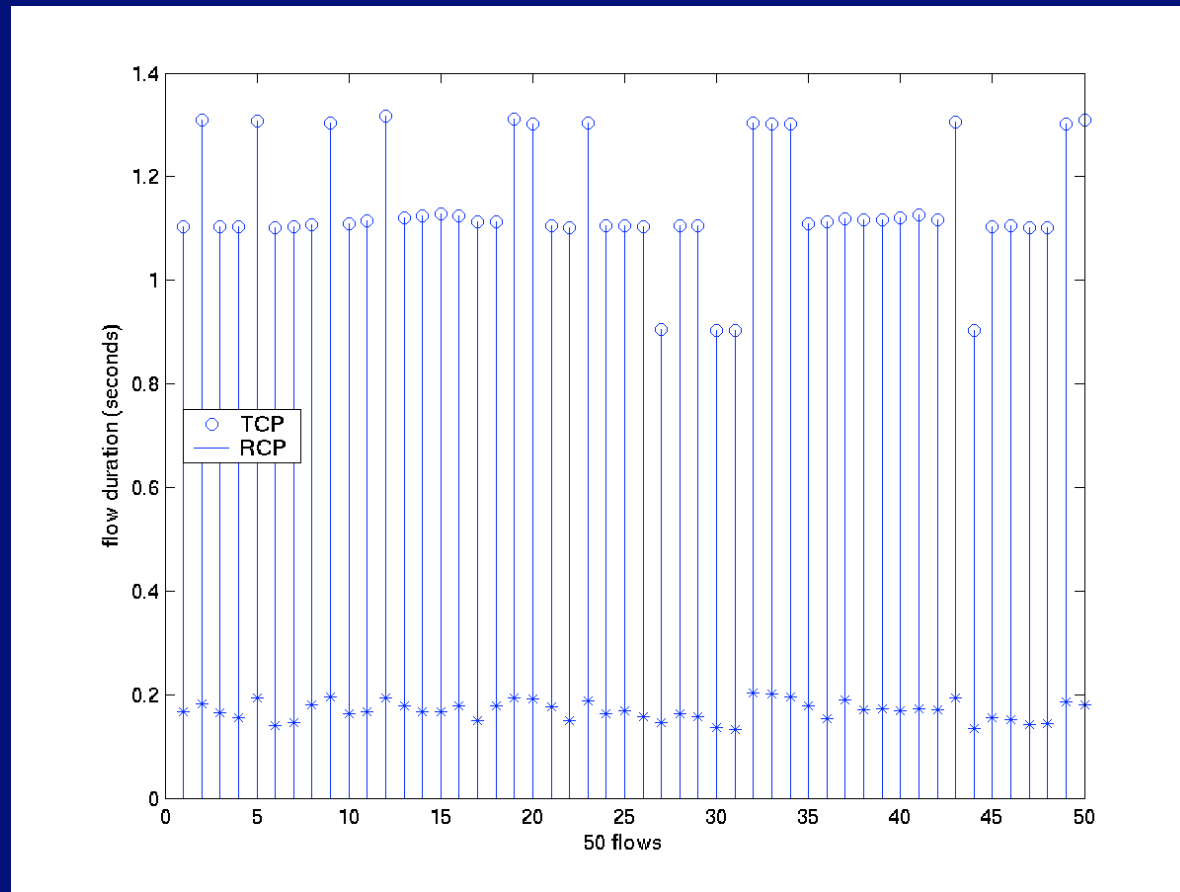
- Router grants  $1/n$  of its capacity to each flow
- This provides fairness amongst flows
- Choosing  $n$  is a tradeoff between
  - Buffer requirements
  - And flow response time

# How to Choose “n”



$\rho$ : average load; Poisscdf(x,  $\rho$ ): cdf of Poisson( $\rho$ ) evaluated at x.

# Simulation Comparison of TCP and RCP



Single bottleneck link of 100Mb/s; RTT=200ms;  $\rho=0.8$ ;  $n=8$ ;  
flow size  $\sim$  Uniform[50, 150]kB

Prob{overload} matches theoretical value in simulation